

Learning with Statistical Queries

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PAC model ✓

What about errors in the labels?

One model for this:

Random Classification Noise Model

each label is flipped with prob. η .

$$(x, l(x)) \quad l(x) = \begin{cases} f(x) & \text{w.p. } 1-\eta \\ 1-f(x) & \text{w.p. } \eta \end{cases}$$

Q. Can we still learn the underlying concept?

E.g. suppose we are learning an OR of Boolean variables.

Let $p_i = \Pr(f(x)=0 \text{ and } x_i=1)$

We can let our hypothesis h be all variables for which $p_i = 0$

(any variable with $p_i = 0$ is in the true OR)

(any variable with $p_i = 0$ or 1 ...
 and no variables for which $p_i > \frac{\epsilon}{n}$.

So can we estimate each $p_i \pm \frac{\epsilon}{2n}$ (say)?

$$p_i = \underbrace{P_{\mathcal{L}}(f(x)=0 / x_i=1)}_{q_i} \cdot P_{\mathcal{L}}(x_i=1)$$

↑ independent of label noise

$$P_{\mathcal{L}_{\text{noise}} \eta}(l(x)=0 / x_i=1) = (1-\eta)q_i + \eta(1-q_i)$$

$$\eta = \eta + q_i(1-2\eta)$$

So if we have LHS, we can subtract η , divide by $(1-2\eta)$.

suffices to approximate to within $\pm \frac{\epsilon}{2n}(1-2\eta)$.
 which we can do from samples!

Statistical Query Model.

Can ask for expectations of bounded functions $\phi(x)$ up to additive

can use \mathcal{Q} functions of $(x, l(x))$ up to additive error.

E.g. $\mathbb{E}(x_i / l(x) = 1)$

$$\mathbb{E}(x_i (1-x_j))$$

$$X: X \times \{0, 1\} \rightarrow [0, 1], \tau > 0$$

↑ ↑
example label

SQ oracle responds with $\mathbb{E}_{\mathcal{D}}(X) \pm \tau$.

X : poly time computable $\tau \geq \frac{1}{\text{poly}}$

Many (almost all) known learning algorithms can be implemented with SQ.

e.g. gradient descent. $L(w) = \mathbb{E}_{x,y}(\ell(x,y,w))$

$$\begin{aligned} \nabla_w L(w) &= \nabla_w \mathbb{E}_{x,y}(\ell(x,y,w)) \\ &= \mathbb{E}_{x,y}(\nabla_w \ell(x,y,w)) \end{aligned}$$

$$\text{SQ.} \rightarrow = E_{x,y} (\nabla_w (L(x,y,w)))$$

Thm. SQ-learnable \Rightarrow PAC learnable
with random classification
noise.

Estimate $P_h(\chi(x, f(x)) = 1)$

Let

$$\underline{\text{CLEAN}} = \{x: \chi(x, 0) = \chi(x, 1)\}$$

$$\underline{\text{NOISY}} = \{x: \chi(x, 0) \neq \chi(x, 1)\}$$

example

$$\chi(x, l) = 1$$

if $x_i = 1$ and $l = 0$.

$$P_h(\chi(x, f(x)) = 1) = P_h(\chi(x, f(x)) = 1 \text{ and } x \in \text{CLEAN}) \\ + P_h(\chi(x, f(x)) = 1 \text{ and } x \in \text{NOISY})$$

Can estimate $P_h(x \in \text{CLEAN})$ — not affected by noise.
and hence first term.

ALSO $P_h(x \in \text{NOISY})$

Also $P_1(x \in \text{NOISY})$

and $P_{1-\eta}(\chi(x, f(x)) = 1 \mid x \in \text{NOISY}) \} q$

$$= (1-\eta)p + \eta(1-p) = \eta + p(1-2\eta)$$

where $p = P_1(\chi(x, f(x)) = 1 \mid x \in \text{NOISY})$

$$\text{Hence } p = \frac{q - \eta}{1 - 2\eta}$$

Need to estimate p to within $\pm \tau$

So need to estimate q to within $\pm \tau(1-2\eta)$

Some functions are hard to learn with SQ.

e.g. parity!

Suppose the hypothesis class is the set of parities. How many SQ(τ) do we need?

The weak Learning PARITY needs $2^{\Omega(n)}$ SQs.

Thm. Weak Learning PARITY needs $\Omega(n)$ queries.

Pf. Recall that $\langle \chi_S, \chi_T \rangle_{\mathcal{D}} = \begin{cases} 1 & S=T \\ 0 & \text{o.w.} \end{cases}$

Take an SQ $\chi(x, \ell(x))$

it is function from $\{-1, 1\}^n \times \{-1, 1\} \rightarrow [-1, 1]$

Let us extend $\{\chi_S\}$ to a basis for $\{-1, 1\}^{n+1}$

$$\chi_S(x, \ell) = \chi_S(x)$$

$$h_S(x, \ell) = \ell(x) \cdot \chi_S(x)$$

These are 2^{n+1} functions.

$$\text{check: } \langle h_S, \chi_T \rangle = 0$$

$$\langle h_S, h_T \rangle = \begin{cases} 1 & S=T \\ 0 & \text{o.w.} \end{cases}$$

So any query g can be written as

So any query g can be written as

$$g(x, l(x)) = \sum_S \alpha_S h_S(x, l) + \underbrace{\sum_S \hat{f}_S \chi_S(x)}_{\text{independent of } l(x)}$$

\therefore

$$\mathbb{E}_D(g(x), l(x)) = \sum_S \alpha_S \mathbb{E}_D(l(x) \chi_S(x)) + g_0$$

Now $l(x) \doteq \chi_T$ for some T . So the first term is 0 except when $S=T$, when we get

$$= \alpha_T$$

T for which $\alpha_T \geq \tau$ is $\leq \frac{1}{\tau^2}$.

(since $\sum_S \alpha_S^2 + \sum_S \hat{f}_S^2 \leq 1$)

So if the SQ oracle responds "0", at most $\frac{1}{\tau^2}$ are eliminated.
... at least $2^n \cdot \tau^2$ queries.

\therefore need τ at least $2 \cdot \tau$ queries.

More generally,

$$\text{SQ-dim}(\mathcal{H}, \gamma) = \max \left\{ d : \begin{array}{l} \exists d \text{ concepts} \\ h_1, h_2, \dots, h_d \in \mathcal{H} \\ \text{s.t. } \|h_i\|^2 \leq 1 \\ \langle h_i, h_j \rangle_D \leq \gamma \end{array} \right\}.$$

Thm. To learn \mathcal{H} in the SQ model, we need $\Omega(d/\gamma)$ queries to $\text{SQ}(\sqrt{\gamma})$.