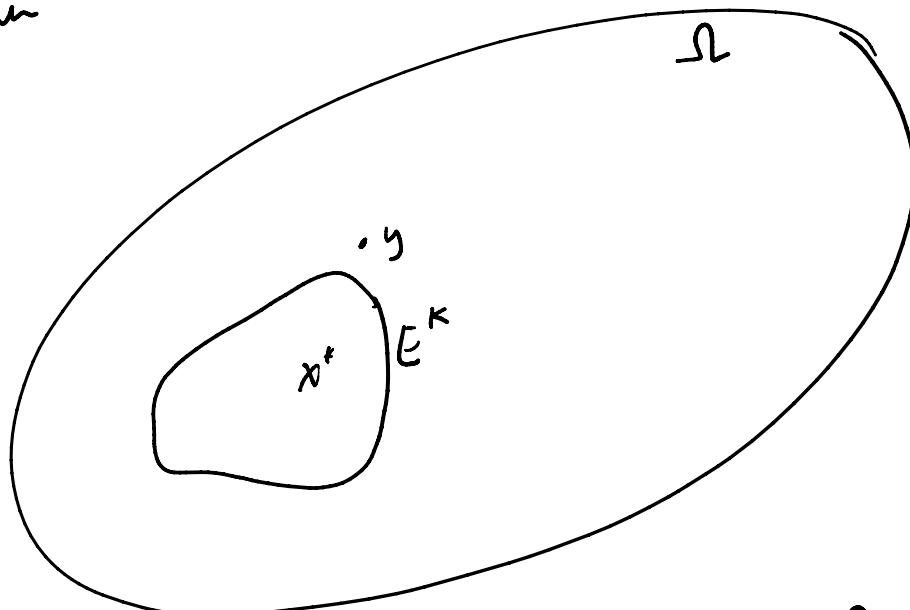


# The Cutting Plane Method

Wednesday, January 15, 2020 6:34 PM

*update*



$$\alpha > \frac{\nu(E_k)}{\nu(\Omega)} . \quad S = (1-\alpha)x^* + \alpha\Omega \quad \nu(S) = \alpha\nu(\Omega)$$

$$y \in S \setminus E_k \Leftrightarrow \nu(S) > \nu(E_k)$$

By separation:  $f(y) \geq f(x^i)$  for some  $i \leq k$ .

$$\exists z \in \Omega : y = (1-\alpha)x^* + \alpha z$$

$$f(x^k) \leq f(y) \leq (1-\alpha)f(x^*) + \alpha f(z)$$

$$f(x^k) - f(x^*) \leq \alpha (f(z) - f(x^*))$$

$$\Rightarrow f(x^k) - f(x^*) \leq \frac{\nu(\Omega)}{\nu(E_k)} (f(x^0) - f(x^*)) \quad (*)$$

Tm. □

Let  $E_0 = \emptyset, E_1, \dots, E_k$  be the sequence of sets  
and  $x^0, \dots, x^k$  be the queries.

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Let  $v: 2^{\mathbb{R}^n} \rightarrow \mathbb{R}_+$  be s.t.

$$(1) \quad v(\alpha E + x) = \alpha v(E)$$

and (2)  $S \subseteq E \Rightarrow v(S) \leq v(E)$ .

then (\*) holds.

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### Cutting plane Method.

Start with  $E_0$ .

Repeat:  $\begin{cases} \text{Choose } x^k \in E_k \\ \text{Using } \nabla f, E_{k+1} \subseteq \left\{ y : \langle \nabla f(x), y - x \rangle \leq 0 \right\} \end{cases}$

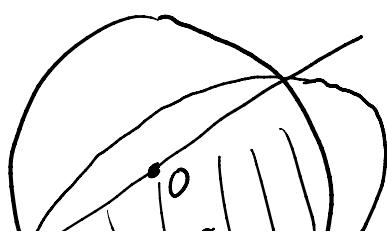
- How to choose  $x^k, E^k$ ?
- How to measure progress?
- Rate of convergence?
- Time to implement each step?

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### Ellipsoid algorithm.

$E_k$  is an ellipsoid.

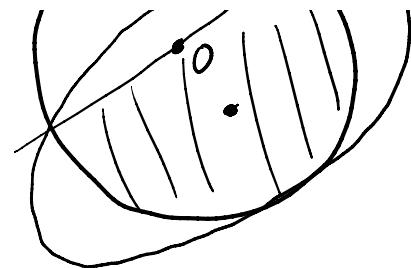
$$E_0 = B(0, R)$$



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$$x^0 = 0$$

$$\nabla f(x_0) (x - x^0) \leq 0$$



$E_1 = \min \text{volume } E$  containing  $E_0 \cap \{x : \nabla f(x^*)^T (x - x^0) \leq 0\}$ .

$x^1 = \text{center of } E_1$ .

Maintain  $x^k$ ,  $E_k = \{x : (x - x^k)^T A_k^{-1} (x - x^k) \leq 1\}$

$$A_0 = R^2 I$$

Repeat :

$$\begin{aligned} x^{k+1} &= x^k - \frac{1}{n+1} \frac{A_k \nabla f(x^k)}{\sqrt{\nabla f(x^k)^T A_k \nabla f(x^k)}} \\ A_{k+1} &= \left( \frac{n^2}{n^2 - 1} \right) \cdot \left( A_k - \frac{2}{n+1} \cdot \frac{A_k \nabla f(x^k) \nabla f(x^k)^T A_k}{\nabla f(x^k)^T A_k \nabla f(x^k)} \right) \end{aligned}$$

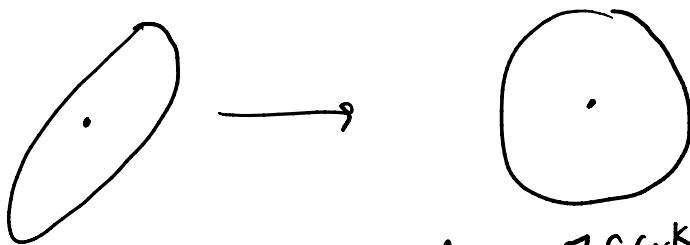
$$\text{Lemma. (1)} \quad \text{vol}(E_{k+1}) \leq e^{-\frac{1}{2n+2}} \cdot \text{vol}(E_k)$$

$$(2) \quad E_k \cap H_k \subseteq E_{k+1}$$

Pf.  $\frac{\text{vol}(E_{k+1})}{\text{vol}(E_k)}$  is maintained by affine transformation.

.. ..  $\sqrt[n+1]{(x - x^k)^T (x - x^k)}$

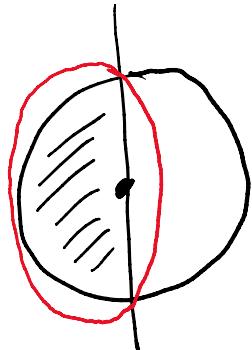
$\overline{\text{Vol}(E_k)}$   
 $E_k = A_k^{\frac{1}{2}} B(0, 1) + x^k$ , so apply  $A_k^{\frac{1}{2}} \cdot (x - x^k)$   
 so that  $\hat{E}_k \rightarrow B(0, 1)$  and  $A_k = I$ .



WLOG assume that  $\nabla f(x^k) = e_1$

$$\text{s.t. } H_k = \{x : e_1^\top (x - o) \leq 0\}$$

i.e.  $x_1 \leq 0$



$$\begin{aligned}
 \text{So } A_{k+1} &= \frac{n^2}{n^2-1} \left( I - \frac{2}{n+1} e_1 e_1^\top \right) \\
 &= \frac{n^2}{n^2-1} \begin{pmatrix} \frac{n-1}{n+1} & & & \\ & 1 & & \\ & & \ddots & \\ & & & 0 \end{pmatrix}
 \end{aligned}$$

$$\left( \frac{\text{Vol}(E_{k+1})}{\text{Vol}(E_k)} \right)^2 = \left| \frac{\det(A_{k+1})}{\det(A_k)} \right|$$

$$= \left( \frac{n^2}{n^2-1} \right)^n \cdot \frac{n-1}{n+1}$$

$$= \left( \frac{n^2}{n^2-1} \right)^{n-1} \cdot \frac{n \cdot n}{(n-1)(n+1)} \cdot \frac{n-1}{n+1}$$

$$= \left( 1 + \frac{1}{n^2-1} \right)^{n-1} \left( 1 - \frac{1}{n+1} \right)^2$$

$$= \left(1 + \frac{1}{n^2-1}\right) \left(1 - \frac{1}{n+1}\right)$$

$$\leq e^{\frac{1}{n^2-1} \cdot (n-1)} - \frac{2}{n+1} = e^{\frac{1}{n+1} - \frac{2}{n+1}} = e^{-\frac{1}{n+1}}$$

$$\text{Vol}(E_{k+1}) \leq e^{\frac{1}{2(n+2)}} \cdot \text{Vol}(E_k).$$


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$$(2) E_k \cap H_k \subseteq E_{k+1}.$$

If  $x \in E_k \cap H_k$ , we have  $\|x\|_2 \leq 1$ ,  $x_1 \leq 0$ .

We need to check:  $(x - x_{k+1})^T \tilde{A}_k^{-1} (x - x_{k+1}) \leq 1$ .

$$x_{k+1} = \left(-\frac{1}{n+1}, 0, \dots, 0\right)^T$$

$$\tilde{A}_{k+1}^{-1} = \frac{n^2-1}{n^2} \cdot \begin{pmatrix} \frac{n+1}{n-1} & 0 & \dots & 0 \\ 0 & 1 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \end{pmatrix}$$

So

$$(x_1 + \frac{1}{n+1}, x_2, \dots, x_n) \cdot \frac{n^2-1}{n^2} \cdot \begin{pmatrix} \frac{n+1}{n-1} & 0 & \dots & 0 \\ 0 & 1 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & 1 \end{pmatrix} \begin{pmatrix} x_1 + \frac{1}{n+1} \\ x_2 \\ \vdots \\ x_n \end{pmatrix}$$

$$= \frac{n^2-1}{n^2} \cdot \left( \frac{n+1}{n-1} \cdot \left(x_1 + \frac{1}{n+1}\right)^2 + x_2^2 + \dots + x_n^2 \right)$$

$$= \frac{n^2-1}{n^2} \cdot \left( \sum x_i^2 + \frac{2}{n-1} x_1^2 + \frac{2x_1}{n-1} + \frac{1}{n^2-1} \right)$$

$$\leq \frac{n^2-1}{n^2} \left(1 + \frac{1}{n^2-1}\right) \leq 1.$$

$$\leq \frac{n-1}{n^2} \left(1 + \frac{1}{n-1}\right) \leq 1.$$

In fact,  $E_{k+1}$  is min volume  $E$  containing  $E_k \cap H_k$ .

Note.  $\nabla f$  is not necessary!

Any  $g$ :  $g^T(x - x^*) \leq 0$  contains  $x^*$  suffices.

i.e. any separating hyperplane of  $\{x : f(x) \leq f(x^*)\}$

App 1. LP.  $\min C^T x, Ax \geq b$ .

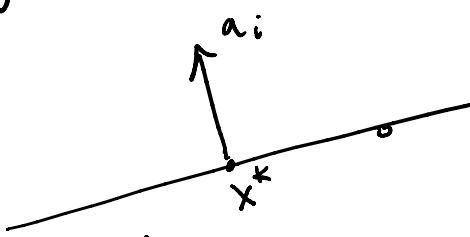
$$\text{Set } f(x) = C^T x + l_{Ax \geq b}(x)$$

$$l(x) = \begin{cases} 0 & Ax \geq b \\ \infty & \text{otherwise.} \end{cases}$$

$$\nabla f(x) = \begin{cases} c & \text{if } Ax \geq b \\ -a_i & \text{if } a_i x < b \quad (\text{if many, pick any one}) \end{cases}$$

$$-a_i^T(x - x^*) \leq 0$$

$$R = \text{Diam } \{x : Ax \geq b\}$$



$$E_0 = B(0, R) \quad V(E) = \text{Vol}(E)^{\frac{1}{n}}, \quad V(\{x : Ax \leq b\}) = r$$

$$E_0 = B(0, R) \quad \mathcal{V}(E) = \text{Vol}(E)^n, \quad \mathcal{V}(\{x : Ax \leq b\}) = r$$

then  $f(x) - f^* \leq \varepsilon (f(x_0) - f^*)$

in at most  $n^2 \log \frac{R}{r\varepsilon}$  steps.

$$\text{Time} = O\left((n^2 + nnz(A)) \cdot n^2 \log \frac{R}{r\varepsilon}\right) \text{ steps.}$$

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