We have studied deterministic finite automata (DFA) and non-deterministic finite automata (NFA). Today we are going to talk about an automata that is the generalization of finite deterministic automata. It is called probabilistic finite automata, or probabilistic FA. It is an automata such that in each state on an input alphabet, has a probability of going to any state.

**Probabilistic FA Example**

**Example 1**

![State Diagram](Image)

The above figure shows state diagram for a simple probabilistic FA (PFA). This PFA has two states $q_0$ and $q_1$. Note that this PFA takes no input. On the contrary, it outputs an alphabet on each transition. Starting from $q_0$, this PFA has a probability of $1/2$ to go to state $q_1$ and output letter $a$, and a probability of $1/2$ to stay in state $q_0$ and output $b$.

This PFA has multiple possible output at each step with different probability. The list of possible outputs is listed in figure 8.1.

<table>
<thead>
<tr>
<th>Output</th>
<th>Probability</th>
<th>Output</th>
<th>Probability</th>
</tr>
</thead>
<tbody>
<tr>
<td>a</td>
<td>$1/2$</td>
<td>ab</td>
<td>$1/3$</td>
</tr>
<tr>
<td>b</td>
<td>$1/2$</td>
<td>aa</td>
<td>$1/6$</td>
</tr>
<tr>
<td></td>
<td></td>
<td>ba</td>
<td>$1/4$</td>
</tr>
<tr>
<td></td>
<td></td>
<td>bb</td>
<td>$1/4$</td>
</tr>
</tbody>
</table>

$t = 1$  
$t = 2$

Figure 8.1: Possible output
Example 2

\[ \begin{array}{ccc}
q_0 & \xrightarrow{1/2, a} & q_1 \\
\xleftarrow{1/2, b} & q_2 & \xleftarrow{2/3, b} \\\n\xleftarrow{2/3, b} & q_1 & \xrightarrow{1/3, a} \\
\end{array} \]

The above shows state diagram for another PFA. This PFA has three state. Similar to the first example, we list all possible output and corresponding probabilities in figure 8.2.

<table>
<thead>
<tr>
<th>Output</th>
<th>Probability</th>
</tr>
</thead>
<tbody>
<tr>
<td>a</td>
<td>1/2</td>
</tr>
<tr>
<td>b</td>
<td>1/2</td>
</tr>
</tbody>
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<td>1/6</td>
</tr>
<tr>
<td>bb</td>
<td>1/6</td>
</tr>
</tbody>
</table>

Figure 8.2: Possible output

**Transition Matrix**

Based on the previous examples, we can see that a probabilistic FA is similar to a usual automata except that in PFA each pair of transition \((q_i, q_j)\) is assigned certain transition probabilities. We define a vector \(\pi^{(t)}\) that describes the distribution of probability over states at the \(t^{th}\) step. The probability of being at state \(q_i\) at \(t^{th}\) step \(P_{\pi}(q^{(t)} = q_i) = \pi^{(t)}_i (0 \leq \pi^{(t)}_i \leq 1, \sum_i \pi^{(t)}_i = 1)\). Next we define a transitive matrix. It is a matrix in which each entry \(p_{i,j}\) represents the probability of going from state \(q_i\) to state \(q_j\). Given the distribution at \(t^{th}\) step \(\pi^{(t)}\), the distribution at \((t + 1)^{th}\) is \(\pi^{(t+1)} = P^T \cdot \pi^{(t)}\).

In the example 2, the transition matrix is

\[
P = \begin{pmatrix}
0 & 1/2 & 1/2 \\
2/3 & 1/3 & 0 \\
0 & 1/3 & 2/3 \\
\end{pmatrix}
\]

Given the initial distribution \(\pi^{(0)} = (1, 0, 0)^T\), the distribution at \(t = 1\) is

\[
\pi^{(1)} = \begin{pmatrix}
0 & 1/2 & 1/2 \\
2/3 & 1/3 & 0 \\
0 & 1/3 & 2/3 \\
\end{pmatrix}^T \cdot \begin{pmatrix} 1 \\ 0 \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \\ 1/2 \\ 1/2 \end{pmatrix}
\]
Similarly the distribution at $t = 2$ is

$$\pi^{(2)} = \begin{pmatrix} 0 & 1/2 & 1/2 \\ 2/3 & 1/3 & 0 \\ 0 & 1/3 & 2/3 \end{pmatrix}^T \begin{pmatrix} 0 \\ 1/2 \\ 1/2 \end{pmatrix} = \begin{pmatrix} 1/3 \\ 1/3 \\ 1/3 \end{pmatrix}$$

When we increase value of $t$, we will see the distribution $\pi^{(t)}$ is getting closer to $(1/4, 3/8, 3/8)^T$.

**Question**: Does $\pi^{(t)}$ converge i.e. $\pi^{(t+1)} = P^T \cdot \pi^{(t)} = \pi^{(t)}(t \to \infty)$?

**Claim**: $\pi^{(t)}$ converges and follows the rule that

$$P^T \cdot \begin{pmatrix} a \\ b \\ b \end{pmatrix} = \begin{pmatrix} a' \\ b' \\ b' \end{pmatrix}$$

**Proof:**

Assume the distribution at $t = k$ is $\pi^{(k)} = (a, b, b)^T$. Given transitive matrix $P$ we have $\pi^{(k+1)} = P^T \cdot \pi^{(t)} = (2/3b, 1/2a+2/3b, 1/2a+2/3b)^T$. To evaluate whether $\pi^{(t)}$ converges towards $u = (1/4, 3/8, 3/8)^T$, we calculate the $\chi^2$ distance between $(u, \pi^{(t)})$ and $(u, \pi^{(t+1)})$. It turns out $\chi^2(u, \pi^{(t+1)}) = 1/9 \chi^2(u, \pi^{(t+1)})$. Clearly as $t$ increases $\pi^{(t)}$ gets closer to $u$, therefore the distribution $\pi^{(t)}$ converges.

**Question**: Does $\pi^{(t)}$ always converge?

**Example 3**

The above figure defines another PFA whose $\pi^{(t)}$ does not converge. All transitions in this PFA have the same probability $1/2$. The distribution for this PFA does not converge since it goes back and forth between $(0, 0, 1/2, 1/2)^T$ or $(1/2, 1/2, 0, 0)^T$.

**Perron–Frobenius Theorem**

**Definition 8.1** A non-negative matrix square $T$ is called **primitive** if $\exists k$ such that $\forall i,j, (T)^k_{i,j} > 0$.

**Definition 8.2** A non-negative matrix square $T$ is called **irreducible** if $\forall i,j, \exists k$ such that $(T)^k_{i,j} > 0$.

**Theorem 8.3** Perron Theorem

For a primitive matrix $P$, $P^k \cdot x \rightarrow v (x \geq 0, x \neq 0)$ and the following conditions hold
1. \( v > 0 \)
2. \( P v = \lambda v \)
3. \( \forall u, P u = \alpha u, \alpha < \lambda \)

**Theorem 8.4 Frobenius Theorem**

For an irreducible matrix \( P \), \( \exists v \) such that the three conditions in Perron theorem hold.

The Perron-Frobenius states that if the transition matrix is primitive, then it has a stationary distribution and will finally converge to it. The stationary distribution is the eigenvector corresponding to the maximum eigenvalue of the transition matrix. But irreducible matrix may not converge to the eigenvector. In fact, a matrix converges if it is irreducible and aperiodic.

**Lemma 8.5** A matrix is primitive \( \iff \) its graph is strongly connected and it is aperiodic.